

CG or PCG method for solving  $Ax=b$ ,  $A$  or  $-A$  is SPD.

PCG: pre conditioning  $Ax=b$

$\text{cond}_2(PA) < \text{cond}(A)$        $PAx=pb$

$\varphi(x) = \frac{1}{2} x^T A x - x^T b$   
 $\nabla \varphi = 0 \implies Ax = b$

CG method:

$x_0, r_0 = b - Ax_0, r_0 \neq 0, p_0 = r_0$

$$\begin{cases} X_{k+1} = X_k + \alpha_k p_k \\ p_k^T r_{k+1} = p_k^T r_k - \alpha_k p_k^T A p_k \\ A p_{k+1} = A r_{k+1} + \beta_{k+1} p_k \end{cases}$$

$\alpha_k = \frac{p_k^T r_k}{p_k^T A p_k}$

$\beta_{k+1} = \frac{-r_{k+1}^T A p_k}{p_k^T A p_k}$

Linear search.  $r_{k+1}^T p_k = 0$

Let  $x_k = x_k - x^*$ ,  $A$  is SPD

$\|e_k\|_2, \|e_k\|_{CG}, \|e_k\|_A = -r$

Not computable

$= \sqrt{e_k^T A e_k}$

$= \sqrt{e_k^T r_k}$

$= \sqrt{(A e_k)^T A e_k}$

$= \|r_k\|_2$

## Convergence

$$r_{k+1} = r_k - \alpha_k A p_k$$

$\alpha_k$  is chosen such that

$\|r_{k+1}\|_{A^T}$  is minimized

$$\|r_{k+1}\|_{A^T} \stackrel{?}{=} \|r_k\|_{A^T} \quad ?$$

$$p_k \in K_R(r_0), \quad r_k \in K_R(r_0)$$

$\{A^T p_j\}_{j=0,1,\dots,k-1}$  is an orthogonal set of  $K_R(r_0)$

$\Rightarrow A$ -conjugate.

Proof:  $p_0 = r_0 = K_0(r_0)$ ,

Assume the conclusions are true for  $j=1, 2, \dots, k$ ,

$$p_{k+1} = r_{k+1} + \beta_{k+1} p_k$$

$$r_{k+1} = r_k - \alpha_k A p_k \in K_k(r_0) = \{r_0, A r_0, \dots, A^{k-1} r_0\}$$

$$r_{k+1} \in K_{k+1}(r_0)$$

$$p_{k+1} = r_{k+1} + \beta_{k+1} p_k \in K_k(r_0) \in K_{k+1}(r_0)$$

We want orthogonal basis

$$K_n(r_0) = \{ r_0, Ar_0, \dots, A^{k-1}r_0 \}$$

$$= \{ r_0, \underline{A^{\perp} p_0}, A^{\perp} p_1, \dots, A^{\perp} p_h \}$$

$\{A^{\perp} p_j\}$  is orthogonal set  $A = UDU^H$   
 $p_i^T A p_j = 0$ , if  $i \neq j$   $A^{\perp} = U D^{\perp} U^H$   
 $= A^{\perp} A^{\perp}$

$$x_{k+1} = x_k + \alpha_k p_k$$

$$b. \quad x_{k+1} = x_{k-1} + \alpha_{k-1} p_{k-1} + \alpha_k p_k$$

$$x_{k+1} = x_0 + \sum_{j=0}^k \alpha_j p_j$$

$$x_{k+1} - x_0 \in K_k(r_0)$$

$$r_{k+1} = r_0 - \sum_{j=0}^k \alpha_j A p_j \in K_{k+1}(r_0)$$

$$e_{k+1} = e_0 + \sum_{j=0}^k \alpha_j p_j$$

$$\|e_{k+1}\|_{A^t} = \min_{\{\alpha_j\}} \left\| e_0 + \sum_{j=0}^k \alpha_j p_j \right\|_{A^t}$$

$$= \min_{\{\alpha_j\}} \sqrt{\left( e_0 + \sum_{j=0}^k \alpha_j p_j \right)^T A \left( e_0 + \sum_{j=0}^k \alpha_j p_j \right)}$$

$$= \min_{\{\alpha_j\}} \left[ e_0^T A e_0 + \sum_{j=0}^k \alpha_j^2 p_j^T A p_j + 2 \sum_{j=0}^k \alpha_j e_0^T A p_j \right]$$

$$\frac{\partial}{\partial \alpha_j} = 2 \alpha_j p_j^T A p_j + 2 e_0^T A p_j = 0$$

$$\alpha_j = \frac{-e_0^T A p_j}{p_j^T A p_j}$$

$$\|e_j\|_{A^t} = \sqrt{e_j^T A e_j}$$

Thm:

$$\|e_k\|_{A^k} \leq 2 \|e_0\|_{A^k}$$

$$\left( \frac{\sqrt{\kappa_2(A)-1}}{\sqrt{\kappa_2(A)+1}} \right)^k$$

Proof sketch:

$$\kappa_2(A) = \text{Cond}_2(A)$$

$$x_k = x_0 + \sum_{j=0}^{k-1} \alpha_j p_j \leq 1$$

$$e_k = e_0 + \sum_{j=0}^{k-1} \alpha_j p_j$$

$$e_k = e_0 - \sum_{j=0}^{k-1} \nu_j A^j r_0$$

$$= e_0 - \sum_{j=0}^{k-1} \nu_j A^{j+1} e_0$$

$$e_k = x^* - x_k$$

$$\rightarrow \kappa_{k-1}(r_0)$$

$$e_1 = e_0 -$$

$$r_0 = b - Ax_0$$

$$= A(A^{-1}b - x_0)$$

$$e_k = \left( I - \sum_{j=0}^{k-1} \nu_j A^{j+1} \right) e_0$$

$$= P_k(A) e_0$$

$$P_k(0) = 1$$

CG and GMRES

A SPD

Any A

$Ax=b$

$$e_k = e_0 - \sum_{j=0}^{k-1} \alpha_j p_j$$

$$= e_0 - \sum_{j=0}^{k-1} \alpha_j A^j r_0$$

$$= e_0 - \sum_{j=0}^{k-1} \alpha_j A^{j+1} e_0$$

$$e_k = A^k b - X_k$$

$$A e_k = r_k$$

$$A x_0$$

$$r_0 = A(A^{-1}b - x_0)$$

$$= A e_0$$

$$= \left( I - \sum_{j=0}^{k-1} \alpha_j A^{j+1} \right) e_0$$

$e_k = p_k(A) e_0$ ,  $p_k(\lambda)$  is a

polynomial of degree  $k$ .  $p_k(0) = 1$

The  $X_k$  from the CG method is the one that minimize  $\|e_k\|_{A^k}$

$$\|e_k\|_{A^k} \leq \frac{\min_{p \in P_k} \|p(A)\|_{A^k}}{p(0)=1} \|e_0\|_{A^k}$$

$$\|A e_k\|_{A^k}$$

$$= \|A\|_2 \|e_k\|_{A^k}$$

$$\sqrt{\lambda_{\min}} \|x\|_2 \leq \|x\|_{A^k} = \sqrt{x^T A x} \leq \sqrt{\lambda_{\max}(A)} \|x\|_2$$

MA

Need  $\gamma > 0$  to find the

minimum.

Chebyshev polynomials

Using Chebyshev approximation

Theorem, we can get

$$\|e_k\|_{A^k} \leq 2 \left( \frac{\sqrt{\kappa_2(A)} - 1}{\sqrt{\kappa_2(A)} + 1} \right)^k \|r_0\|_{A^k}$$

A is not sym. SPD

M.I.  $A^T A x = A^T b$

$$\text{Cond}(A^T A) = (\text{Cond}(A))^2$$

GMRES:

$$x_k = x_0 + \sum_{j=0}^{k-1} \gamma_j A^j r_0$$

$$r_k = r_0 - \sum_{j=0}^{k-1} \gamma_j A^{j+1} r_0$$

Can not use  $\|e_k\|_{A^k}$ , Does not exist.

$$\|r_k\|_2 = \min_{P_k \in P_k, p(0)=1} \|p(A)\|_2 \|r_0\|_2$$

Need to find an orthogonal basis in  $K_k(r_0)$ . Only add one vector at one iteration.



Arnoldi process.

$$r_0, Ar_0, \dots, A^{k-1}r_0$$

$$v_1 = r_0 / \|r_0\|_2$$

$$h_{21}, v_2 = Av_1 - h_{11}v_1$$

$$\left( \begin{array}{l} 0 = v_1^T Av_1 - h_{11} \end{array} \right.$$

$$v_2^T v_1 = 0$$

$$v_2^T v_2 = 1$$

$$h_{11} = v_1^T Av_1$$

$$h_{21} = \|Av_1 - h_{11}v_1\|_2$$

Assume we have  $v_1, v_2, \dots, v_k$

$$h_{k+1,k}, v_{k+1} = Av_k - \sum_{j=1}^k h_{jk} v_j \quad v_i^T v_j = \begin{cases} 1 & i=j \\ 0 & \end{cases}$$

$$h_{jk} = v_j^T Av_k, \quad h_{k+1,k} = \|Av_k - \sum_{j=1}^k h_{jk} v_j\|_2$$

$$A [v_1 \ v_2 \ \dots \ v_k] = [v_1 \ v_2 \ \dots \ v_{k+1}]_{n \times (k+1)}$$

$$AV_k = V_{k+1}H_k$$

GMRFS

$$x_k = x_0 + \sum_{j=1}^k \beta_j v_j$$

$$\begin{bmatrix} h_{11} & h_{12} & \dots & h_{1k} \\ h_{21} & h_{22} & \dots & h_{2k} \\ \vdots & \vdots & \ddots & \vdots \\ 0 & & & h_{kk} \\ & & & h_{k+1,k} \end{bmatrix}_{(k+1) \times k}$$

$$x_k = x_0 + v_k y_k$$

multiply by A

$$\|v_k\|_2 = \|r_0 - v_{k+1} H_k y_k\|_2$$

The GMRFS method, we choose

$y_k$  to minimize  $\|v_k\|_2$

$$= \|v_{k+1}^T (v_{k+1}^T r_0 - H_k y_k)\|_2$$

$$= \| \|r_0\|_2 e_1 - H_k y_k \|_2$$

$$\begin{bmatrix} x & x & x & \dots & x \\ x & x & x & x & x \\ \vdots & & & & \\ 0 & & & x & x \\ & & & x & x \end{bmatrix} \begin{bmatrix} y_1 \\ \vdots \\ y_k \end{bmatrix} = \begin{bmatrix} \|r_0\|_2 \\ 0 \\ \vdots \\ 0 \end{bmatrix}$$

least squares solution,  $\rightarrow QR$ .

Given's rotation