

Summarize convergence of J, G-S,

General matrices:

SOR(ω)

SCDD J, G-S

Weakly COD, +, irreducible G-S, J.

A: SPD, G-S, SOR(ω), $0 < \omega < 2$

→ Poisson eqn. (1D, 2D) $D = \sigma I$ $A = D - L - U$

$$\rho(R_J) = \max_{1 \leq i \leq n} \left| 1 + \frac{\lambda_i(A)}{2} \right|$$

$$\det(\lambda I - D^{-1}(L+U)) = 0 \rightarrow A$$

$D = \sigma I$

$$\det(\lambda D - (L+U)) = 0 \quad \hat{A} = \begin{bmatrix} \lambda^2 & & \\ & \ddots & \\ & & \lambda^2 \end{bmatrix}$$

$$\det(\underline{2(\lambda-1)I} - A) = 0$$

$$\lambda_i'(A) = 2(\lambda_i(R) - 1)$$

$$\det(\lambda' I - A) = 0$$

$$\lambda_i(R) = 1 + \frac{\lambda_i(A)}{2}$$

Lemma: Eigenvalues of

$$A = \begin{bmatrix} \alpha & \beta & & & 0 \\ \beta & \alpha & & & \\ & & \ddots & & \\ & & & \ddots & \\ 0 & & & & \beta & \alpha \end{bmatrix}_{n \times n}$$

are

$$\lambda_k = \alpha + 2\beta \cos \frac{\pi k}{n+1}, \quad k=1, 2, \dots, n.$$

$$x_{k,j} = \sin \frac{k\pi j}{n+1}, \quad j=1, 2, \dots, n.$$

check

$$A X_k = \lambda_k X_k$$

$$-(\beta u')' + \alpha u = \lambda u, \quad 0 < x < 1$$

$$u(0) = 0, \quad u(1) = 0$$

$$A = \begin{bmatrix} -2 & & & 0 \\ & 1 & & \\ & & \ddots & \\ 0 & & & 1 \\ & & & & -2 \end{bmatrix}$$

$$\lambda_k = -2 + 2 \cos \frac{\pi k}{n+1}$$

$$Ax = \lambda x$$

$$A^T A = A^2 x = \lambda^2 x$$

$$\text{Cond}_2(A) = K(A) = \frac{\max |\lambda_i(A)|}{\min |\lambda_i(A)|} \approx 4$$

$$\lambda_k = -4 \left(1 - \cos \frac{\pi k}{n+1} \right)$$

$$= -4 \sin^2 \frac{\pi k}{2(n+1)}$$

$$\approx -4 \left(\frac{\pi k}{2(n+1)} \right)^2$$

$$\text{Cond}_2(A) = \frac{4(n+1)^2}{4(\pi k)^2}$$

$$\sim (n+1)^2$$

$$\rho(R_J) = \max_{1 \leq k \leq n} \left| 1 + \frac{\lambda_k(A)}{2} \right| \quad \left| 1 - \frac{\lambda_k(A)}{2} \right|$$

$$= \max_{1 \leq k \leq n} \left| 1 - \frac{2(1 - \cos \frac{\pi k}{n+1})}{2} \right|$$

$$= \max_{1 \leq k \leq n} \left| \cos \frac{\pi k}{n+1} \right| < 1 \quad k=1, 2, \dots$$

$$\cos \frac{\pi k}{n+1} = 1 - \left(\frac{\pi k}{n+1} \right)^2$$

The Jacobi iterative method needs $O(n^2)$ for convergence. It's slow, but simple!

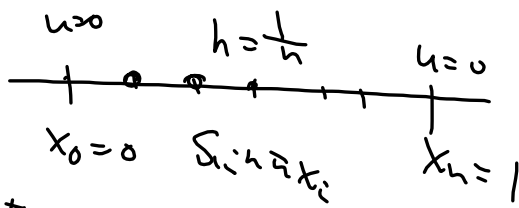
$$-u'' = \lambda u$$

$$u(0) = u(1) = 0$$

$u=0$, trivial soln.

S-L eigenvalue prob.

✓ $\lambda_n = (n\pi)^2, u(x) = \sin n\pi x, n=1, 2, \dots$



$$\sin n\pi x_i$$

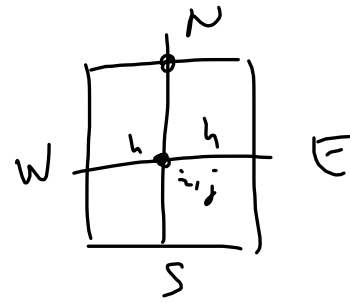
$$\sin \frac{\pi x_i}{h}$$

$$\lambda_k = -2 + 2 \cos \frac{\pi k}{n}$$

$k=1, 2, \dots, n-1$

$$x_{i,j} = \sin \frac{\pi k_j}{h}, \quad j=1, \dots, n-1$$

2D: $u_{xx} + u_{yy} = f(x, y)$
 $= \lambda u$



$$\frac{u_{i-1,j} + u_{i+1,j} + u_{i,j-1} + u_{i,j+1} - 4u_{i,j}}{h^2} = f_{i,j}$$

$i, j = 1, 2, \dots, n-1$

$$u_{i,j} = \frac{1}{4} \left(\dots \right)$$

$$\lambda_{i,j} = - \left(4 - 2 \left(\cos \frac{\pi i}{n} + \cos \frac{\pi j}{n} \right) \right)$$

$i, j = 1, \dots, n-1$

$$\rho(R_5) = \max_{1 \leq i, j \leq n-1} \left| 1 + \frac{\lambda_{i,j}(A)}{4} \right|$$

$$= \max_{1 \leq i, j \leq n-1} \left| 1 - \frac{1}{4} \left(4 - 2 \left(\cos \frac{\pi i}{n} + \cos \frac{\pi j}{n} \right) \right) \right|$$

$$= \max_{1 \leq i, j \leq n-1} \frac{1}{2} \left| \cos \frac{\pi i}{n} + \cos \frac{\pi j}{n} \right|$$

$$\approx \cos \frac{\pi}{n} \sim 1 - O\left(\frac{1}{n^2}\right)$$

$$\rho(R_J) = \max \left| 1 + \frac{\lambda_{i,j}(A)}{2} \right| \quad 1D$$

$$\left| 1 + \frac{\lambda_{i,j}(A)}{4} \right| \quad 2D$$

The Best ω : $SOR(\omega)$, $0 < \omega < 2$.

This is a necessary condition for $SOR(\omega)$ to converge.

$$\omega = \frac{-\rho_1}{2, 1} \rightarrow \text{diverge!}$$

For general matrix A , we don't know the best ω . For Poisson Eqn. we know the

$$\text{best } \omega \quad 0 < \omega_{\text{opt}} = \frac{2}{1 + \sqrt{1 - \rho^2(K_J)}} < 2$$

Poisson Eqn. 1D $\frac{U_{i-1} - 2U_i + U_{i+1}}{h^2} = f_i$

2D $\frac{U_{i-1,j} + \dots - 4U_{ij}}{h^2} = f_{ij}$ $i=1, 2, \dots, n-1$ $\begin{cases} U_0 = \alpha \\ U_n = \beta \end{cases}$ $h \rightarrow \frac{1}{n+1}$

$D = -2I$, $D = -4I$ $-A$ is SPD

$\rho(R_J) = \max \left| 1 + \frac{\lambda_L(A)}{2} \right| \left| \frac{\lambda_L(A)}{4} \right| < 1$

The No. of iterations for Gauss-J is $O(N^2)$

Thm. For 1D 2D Poisson eqns

SOR(ω_{opt}) is $O(N)$

(A): $(\lambda(\text{SOR}(\omega)) + \omega - 1)^2 = \lambda(\text{SOR}(\omega)) \cdot \omega^2 \lambda_J^2$

(B): $\omega_{opt} = \frac{2}{1 + \sqrt{1 - \rho^2(R_J)}}$ $D = -4I$

Key: $R_{SOR} = (D - \omega L)^{-1} ((1 - \omega)D + \omega U)$

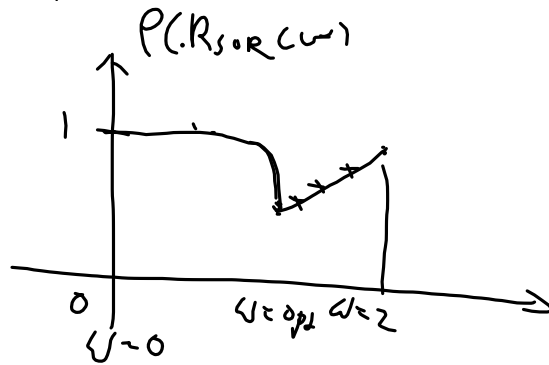
FFT $O(n^2 \log n)$
 ~~$O(n^2)$~~
 $\sim O(n^2)$
 Multigrid

$P(R_{SOR})$ We would choose ^{rather} ω larger than smaller. $0 < \omega < 2$ for any A.

$$P(R_{SOR}) = \begin{cases} 1 - \omega + \frac{1}{2} \omega^2 P(R_J)^2 + \omega P(R_J) \\ \sqrt{1 - \omega + \frac{\omega^2}{4} P(R_J)^2} \end{cases}$$

$P(R_{SOR}) < 1$ if $0 < \omega < 2$.

$\omega < 1$ $0 < \omega < \omega_{opt}$



$$\omega_{opt} = \frac{2}{1 + \sqrt{1 - \frac{P(R_J)^2}{n}}} = \frac{2}{1 + \frac{\pi}{n+1}}$$

Other extensions: JOR(ω)

$$x^{k+1} = (1 - \omega)x^k + \omega \bar{x}_J^{k+1}$$

$$i = 1, 2, \dots, n \rightarrow \text{SOR}(\omega)$$

Then $i = n, n-1, \dots, 1$.

SSOR symmetric SOR(ω). alternatively.

Algebraic eigenvalue problems.

Given $A \in \mathbb{R}^{n \times n}$, if there is $\lambda \in \mathbb{C}$,

and $x \in \mathbb{R}^n$, $x \neq 0$, such that

$$Ax = \lambda x, \quad (\lambda, x) \text{ is eigen-pair of } A.$$

λ can be found theoretically by solving

$$\det(\lambda I - A) = \sum_{k=0}^n \binom{n}{k} \lambda^{n-k} (-1)^k \text{Trace of } A^k \neq 0$$

The characteristic polynomial of A is $P_A(\lambda) = \det(\lambda I - A) = 0$

There are

$$\lambda_1, \lambda_2, \dots, \lambda_n$$

$$P_A(x) = 0, \text{ roots } f(x) = 0 \text{ zero}$$

can be multiple eigenvalues of A .

$$\lambda_1 = 1 - 2i$$

$$\lambda_2 = 1 + 2i$$

Complex roots are in pair

If λ is an eigenvalue of A , what are eigenvalues of $A^2, \dots, A^k, \dots, A^{-k}$

$$Ax = \lambda x, \quad A^2 x = \lambda Ax = \lambda \lambda x = \lambda^2 x \quad \frac{1}{\lambda^k}$$

$$\begin{aligned} \text{Cond}(A^T A) &= \text{Cond}(A^2) = \|A^2\|_2 \| (A^2)^{-1} \|_2 \\ &= \lambda_{\max}(|\lambda_i^2(A)|) / \min(|\lambda_i^2(A)|) \end{aligned}$$

$$A^{-1}, (A^{-1})^2$$

$$Ax = \lambda x \quad x = \lambda A^{-1} x$$

$$\frac{1}{\lambda} x = A^{-1} x$$

If $\lambda = 0$, A is singular, Can find non-zero $x \neq 0$, $Ax = 0$.

$$\text{Cond}_2(A) = K(A) = \frac{\max(|\lambda_i(A)|)}{\min(|\lambda_i(A)|)}$$

$$A = \begin{bmatrix} 2 & 0 & 0 \\ 0 & 2 & -1 \\ 0 & -1 & 2 \end{bmatrix}$$

$$\underline{\lambda_1 = 2}$$

$$A = \begin{bmatrix} D_1 & \\ & D_2 \end{bmatrix}$$

$$\det \begin{pmatrix} \lambda - 2 & 1 \\ 1 & \lambda - 2 \end{pmatrix} = 0$$

$$\lambda_2(A) = \{ \lambda_2(D_1) \} \cup \{ \lambda_2(D_2) \}$$

$A \in \mathbb{R}^{n \times n}$. Assume A has a complete
eigen system, or A is diagonalizable

$$v_1, v_2, \dots, v_n$$

$$A v_i = \lambda_i v_i, \quad i = 1, 2, \dots, n \quad \|v_i\| = 1$$

$$Q^{-1} A Q = D = \begin{pmatrix} \lambda_1 & & 0 \\ & \lambda_2 & \\ 0 & \ddots & \\ & & \lambda_n \end{pmatrix}$$

$$Q^T A Q$$

$$\frac{|\lambda_1|}{= P(A)} > |\lambda_2| \geq |\lambda_3| \geq \dots \geq |\lambda_n|$$

essential

Observation: $x^0 \neq 0, \|x^0\|=1,$

$$x^1 = Ax^0, \quad x^2 = Ax^1, \quad \dots \quad x^{k+1} = Ax^k \\ = \dots = A^{k+1} x^0$$

$$x^0 = \sum_{j=1}^n \alpha_j v_j$$

$$x^1 = Ax^0 = \sum_{j=1}^n \alpha_j A v_j = \sum_{j=1}^n \alpha_j \lambda_j v_j$$

$$x^2 = \sum_{j=1}^n \alpha_j \lambda_j^2 v_j$$

...

$$x^{k+1} = \sum_{j=1}^n \alpha_j \lambda_j^{k+1} v_j \quad \left| \frac{\lambda_i}{\lambda_1} \right| < 1, \quad i=2, \dots, n$$

$$= \lambda_1^{k+1} \left(\alpha_1 v_1 + \alpha_2 \left(\frac{\lambda_2}{\lambda_1} \right)^{k+1} v_2 + \dots \right)$$

$$x^{k+1} \approx \lambda_1^{k+1} \alpha_1 v_1 + \alpha_n \left(\frac{\lambda_n}{\lambda_1} \right)^{k+1} v_n$$

$$Ax^{k+1} = \lambda_1^{k+2} \alpha_1 v_1 \rightarrow \infty \text{ if } |\lambda_1| \geq 1$$

$$\rightarrow 0 \text{ if } |\lambda_1| < 1.$$

Solu: Scaling.

Start from x^0 , $\|x^0\|_2 = 1$, $\begin{bmatrix} 1 \\ -1 \end{bmatrix} / \sqrt{2}$

$$\begin{cases} y_1 = Ax_0 \\ x_1 = y_1 / \|y_1\|_2 \\ \mu_1 = x_1^T A x_1 \end{cases}$$

1 mu

Algorithm

Given $x^0 \neq 0$,

for $i=1$, until converges.

$$y_{i+1} = A x_i$$

$$x_{i+1} = y_{i+1} / \|y_{i+1}\|_2$$

$$\mu_{i+1} = x_{i+1}^T A x_{i+1}$$

end

$$\begin{cases} y_2 = A x_1 \\ x_2 = y_2 / \|y_2\|_2 \\ \mu_2 = x_2^T A x_2 \end{cases}$$

Raileigh quotient

$$\frac{x^T A x}{x^T x}$$

Theorem: (i) $\lim_{k \rightarrow \infty} \mu_k = \lambda_1$, (ii) $\lim_{k \rightarrow \infty} X_k = v_1$

$$A v_1 = \lambda_1 v_1$$

Proof:

$$X_{k+1} = \frac{y_{k+1}}{\|y_{k+1}\|_2} = \alpha A X_k = \alpha A^2 X_{k-1}$$

$$= \alpha A^{k+1} x_0 \rightarrow 0$$

$$= \alpha \lambda_1^{k+1} \left(\alpha_1 v_1 + \alpha_2 \left(\frac{\lambda_2}{\lambda_1}\right)^{k+1} v_2 + \dots + \alpha_n \left(\frac{\lambda_n}{\lambda_1}\right)^{k+1} v_n \right) \rightarrow 0$$

$$\lim_{k \rightarrow \infty} X_{k+1} = \beta v_1$$

Since $\|v_i\|_2 = 1$

$$\|X_{k+1}\|_2 = 1, \quad \beta = 1.$$

$$X^T X = \|X\|_2^2$$

$$\lim_{k \rightarrow \infty} \mu_k = \lim_{k \rightarrow \infty} X_k^T A X_k = \lambda_1 v_1^T v_1 = \lambda_1$$

Real condition: $|\lambda_1| > |\lambda_2|$

The theorem is still valid!

Complexity: How many matrix-vector multiplication. $y_{i+1} = A y_i \quad O(n^2)$

$$y_i^T A x_i \quad O(n)$$

How do we find the smallest magnitude of all eigenvalues.

$$A \quad \lambda_1, \lambda_2, \dots, \lambda_n \quad Av_n = \lambda_n v_n.$$

$$\underline{|\lambda_n|} < |\lambda_{n-1}| \leq \dots \leq |\lambda_1|$$

A^{-1} : The eigenvalues are $\det(A) \neq 0$

$$\frac{1}{\lambda_n} \quad \frac{1}{\lambda_{n-1}} \quad \dots \quad \frac{1}{\lambda_1}$$

$$\left| \frac{1}{\lambda_n} \right| > \frac{1}{|\lambda_{n-1}|} \geq \frac{1}{|\lambda_{n-2}|} \geq \dots \geq \frac{1}{|\lambda_1|}$$

$$\begin{cases} y_{i+1} = A^{-1} x_i & \Rightarrow A y_{i+1} = x_i \\ x_{i+1} = y_{i+1} / \|y_{i+1}\|_2 \\ \mu_{i+1} = x_{i+1}^T A^{-1} x_{i+1} \end{cases}$$

The inverse power method $PA=LU$

$$A y_{i+1} = x_i$$

$$\begin{cases} L U y_{i+1} = P x_i \\ x_{i+1} = y_{i+1} / \|y_{i+1}\|_2 \\ \mu_{i+1} = x_{i+1}^T \underline{A} x_{i+1} \end{cases}$$